Implementing Counters with Decay

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Decaying Counters

- Counters for counting events used in search heuristics
 - AFC (accumulated failure count = weighted degree)
 count how often a constraint has failed
 sum up over all constraints of a variable
 [Boussemart et al., ECAI 2004]
 - activity
 count how often a variable has been pruned
 [Michel, Van Hentenryck, CPAIOR 2012]
- Values should decay as search progresses
 - random restarts: must decay
- Decay: if counter not incremented, scale with decay factor γ (0 < γ ≤ 1)

Problem

- Data structures
 - counters $c_1, ..., c_n$ (floating point numbers, C++: **double**)
- Implementation

```
proc inc(i) = c_i \leftarrow c_i + 1; forall j \neq i do c_j \leftarrow c_j \cdot \gamma;
fun val(i) = return c_i;
• complexity of inc(i): O(n)
```

- But: counter set possibly not known when doing inc(i)
 - no **forall** possible
- But: inc(i) might suffer from contention with parallel search
 - all counters must be locked and O(n) operations!

Solution

- Data structures
 - counters $c_i = \langle n_i, t_i \rangle$ " $\langle value, timestamp \rangle$ "
 - global timestamp *t*
- Implementation

proc inc(i) =
$$n_i \leftarrow n_i \cdot \text{pow}(\gamma, t - t_i) + 1$$
; $t \leftarrow t + 1$; $t_i \leftarrow t$; **fun** val(i) = $n_i \leftarrow n_i \cdot \text{pow}(\gamma, t - t_i)$; $t_i \leftarrow t$; **return** n_i ;

- complexity of inc(*i*): O(1)
- AFC: often $O(n^2)$ calls to val(i) for each call to inc(i)
- Optimizations
 - cache for likely exponents of pow (which is expensive to compute)
 - do not do anything for $\gamma=1$